15-150 Principles of Functional Programming Lecture 13 February 27, 2025 Michael Erdmann **Exceptions**, n-Queens,

& more success continuations

Exceptions

### Declaring Raising Handling

Exceptions are useful for signaling errors, including violations of invariants. Exceptions can also be useful for control flow, analogous to continuations.



This line of code declares a new exception (constructor) called Silly. (If Silly already exists, this new Silly will

(If Silly already exists, this new Silly will shadow the earlier Silly.)



Silly : exn 1 the type for exceptions raise Silly: 'a if 3=4 then raise Silly else 9 : int So instantiate 'a to be int

and the expression reduces to value 9.

What about the type & value of

## if 3=3 then raise Silly else 0?

The type is int. The expression does <u>not</u> reduce to a value. Instead SML will print <u>uncaught exception Silly</u>. We will discuss handling exceptions shortly.



<u>Handling</u> General form of an exception handler for expressione: <u>e handle</u>  $P_1 \Rightarrow e_1$   $|P_2 \Rightarrow e_2$  $|Pn \Rightarrow e_n$ 

e, e, ..., en are expressions; must have the same type. Pis..., Pn are patterns; must match exceptions. If e reduces to a value V, that value is returned. If e instead raises an uncaught exception E, the handler will try to match E against the patterns Pis..., Pn (in sequential order). If E matches Pi, then SML will evaluate e;. If no pattern matches, E remains uncaught.

## $\frac{f_{un}}{f(x,o)} = \frac{raise}{raise} Rdiv (x * x)$ $\int f(x,n) = \frac{if}{n < 0} \frac{then raise}{raise} Silly$ $\frac{else}{x} (real n)$

 $\frac{fun}{g(x,n)} = f(x,n) \frac{handle}{Rdiv(v)} \Rightarrow 0.0$   $|Rdiv(v) \Rightarrow v$ 

# $\frac{f_{un}}{f(x,o)} = \frac{raise}{raise} Rdiv (x * x)$ $\int f(x,n) = \frac{if}{n \times 0} \frac{raise}{raise} Silly$ $\frac{else}{x} (real n)$

 $\frac{fun}{g(x,n)} = f(x,n) \frac{handle}{Rdiv(v)} \Rightarrow 0.0$   $|Rdiv(v) \Rightarrow v$ 

what are the values of:

- g(3.0,0) (3.0, 9.0
- g (3.0, 2) 1.5
- g (3.0, ~1) ~ 0.0

# $\frac{f_{un}}{f(x,o)} = \frac{raise}{raise} Rdiv (x * x)$ $\int f(x,n) = \frac{if}{n \times 0} \frac{then raise}{raise} Silly$ $\frac{else}{x} (real n)$

$$fun g(x,n) = f(x,n) \underline{handle Silly} \Rightarrow 0.0$$
  
Suppose g does not  
handle Rdiv.  
$$what now are the values of:$$
$$g(3.0, 0) no value; uncaught exception Rdiv$$

$$g(3.0, \sim 1) \longrightarrow 0.0$$

n-Queens

Place n queens on a square nxn board without any two queens threatening each other.

(A queen threatens all locations in the same column, in the same row, and on lines with slope ±1 that pass through the queen's position.)

### Three Implementations

- Exceptions
- Options
  - · Continuations

We will use these programming styles for Search

#### Example: Place 4 queens on a 4x4 board:

#### Start by placing a queen in column 1 and row 1:



























### column 3





column 3











#### **OH NO!** We cannot place the third queen!



## column 3

#### **OH NO! We cannot place the third queen!** Let's backtrack to the placement of second queen.





#### Let's backtrack to the placement of second queen.



#### Let's try a new placement for the second queen.



#### Let's try a new placement for the second queen.



column 2










column 3



## That succeeds!

column 3



























## **OH NO!** We cannot place the fourth queen!





## **OH NO! We cannot place the fourth queen!** Let's backtrack to the placement of third queen.





#### Let's backtrack to the placement of third queen.











## **OH NO!** Again cannot place the third queen!



# column 3

## **OH NO!** Again cannot place the third queen! Again backtrack to the placement of second queen.





## Again backtrack to the placement of second queen.







## **OH NO!** We cannot place the second queen!



column 2

## **OH NO! We cannot place the second queen!** Let's backtrack to the placement of *first* queen.



column 1

#### Let's backtrack to the placement of *first* queen.









### Eventually, placement of second queen succeeds:



#### Then placement of third queen succeeds:



row 1

column 3





### Eventually, placement of **fourth** queen **succeeds**:



column 4

## Solution obtained:



## Code Overview

- (i,j) refers to ith column & jth row.
- Try to add a queen to column i, given threat-free queen placements in columns 1,...,i-1.
- Try successive rows, i.e., positions (i, 1), (i, 2)...
- If position (i,j) is threat free, place
   i<sup>th</sup> gueen there and move on to column i+1.
- If no position is threat-free in column i, backtrack to column i-1, undo the prior placement of a gueen in that column and search for a new placement.

(\* threat : int \* int → int \* int → bool
 Decide whether two queen positions
 threaten each other.
\*)
fun threat (x,y) (a,b) =
 x=a orelse y=b orelse x+y=a+b orelse x-y=a-b

(\* threat : int \* int -> int \* int -> bool \*) <u>fun</u> threat (x,y) (a,b) = x=a <u>orelse</u> y=b <u>orelse</u> x+y=a+b <u>orelse</u> x-y=a-b

(\* conflict : int int -> (int + int) list -> bool Decide whether a given queen position is threatened by any queen position in a list of queen positions. fun conflict p = List. exists (threat p) List. exists : ('a -> bool) -> 'a list -> bool

(\* addqueen : int \* int \* (int \* int) list → (int \* int) list try : int → (int \* int) list queens : int → (int \* int) list

- addqueen (i,n,Q) tries to place all remaining queens on an n×n board, starting in column i, assuming Q describes conflict-free queen placements in columns 1,...,i-1.
- addqueen uses local helper function try.
   try (j) starts its search from position (i,j).
- · queens (n) tries to place all queens on an nxn board.
- These functions raise exception Conflict when unsuccess ful.
   \*)

exception Conflict <u>fun</u> addqueen (i, n, Q) = <u>let</u> fun try j = (if conflict (i, j) Q then raise Conflict else if i=n then (i,j) :: Q <u>else</u> addqueen (i+1, n, (i, j):: Q)handle Conflict => if j=n then raise Conflict else try (j+1) in try 1 end fun queens n = addqueen (1, n, [])
queens  $4 \hookrightarrow [(4,3), (3,1), (2,4), (1,2]$ queens  $1 \hookrightarrow [(1,1)]$ queens 2 does not return a value. Instead, exception Conflict

is uncaught at top level.

Implementation using options

$$\frac{fun}{fun} addqueen (i, n, Q) = \frac{let}{fun} try j = (\underline{case} (\underline{if} conflict} (i, j) Q then NONE) 
else if i = n then SOME((i, j)::Q) 
else addqueen (i + 1, n, (i, j)::Q)) 
of NONE  $\Rightarrow \underline{if} j = n \underline{then} NONE$    
else try (j + 1)   
l result  $\Rightarrow$  result)   
in try 1   
end   
fun queens  $n = addqueen (1, n, [])$$$

## Implementation using continuations

(\* addqueen : int \* int \* (int \* int) list  $\rightarrow ((int * int) | ist \rightarrow 'a)$   $\rightarrow (unit \rightarrow 'a)$ 'a try : int -> 'a queens : int -> (int \* int) list option (\* Here we have the top-level queens function **\***) again return a list option of queen placements. \*) <u>fun</u> addqueen (i,n,Q) sc fc = <u>let</u><u>fun</u> try j =  $\frac{\text{let}}{\text{fun}} \text{fc}'() =$ if j=n then fc() else try(j+1) in if conflict (i,j) Q then fc'() else if i=n then sc((i,j)::Q)else addqueen (i+1,n,(i,j)::Q) sc fc' end In try 1 end <u>fun</u> queens n = addqueen (1, n, []) Some  $(fn() \Rightarrow NONE)$  More powerful continuations

We will allow success continuations to take failure continuations as arguments.

Doing so increases expressive power. We can then solve more problems simply by changing continuations slightly.

<u>datatype</u> 'a tree = Empty | Node of a tree \* 'a \* 'a tree (\* find : ('a → bool) → 'a tree  $\rightarrow$  ('a  $\rightarrow$  (unit  $\rightarrow$ 'b)  $\rightarrow$ 'b)  $\rightarrow$  (unit  $\rightarrow$  'b) ~ ~ **\***) <u>fun</u> find p Empty sc fc = fc() | find p (Node(l,x,r)) sc fc = let fun fonew () = find pl sc (fn () => find pr sc fc) in if p(x) then sc x fenew <u>else</u> fenew() The success continuation receives element **x** and the failure continuation (fcnew end says what to do if p(x) had been false).

fun even  $n = (n \mod 2 = 0)$ 

(\* find first even integer encountered in pre-order traversal.) fun findfirst T = find even T  $(\underline{fn} x \Rightarrow \underline{fn} f \Rightarrow SOME(x))$  $(f_n() \Rightarrow NONE)$ (\* accumulate list of all even integers. \*) fun findall T = find even T  $(\underline{fn} \times \Rightarrow \underline{fn} f \Rightarrow \times :: f())$  $(\underline{fn}() \Rightarrow [1])$ (\* count all the even integers. \*) fun count T = find even T  $(f_n x \Rightarrow f_n f \Rightarrow 1 + f())$  $(f_n() \Rightarrow 0)$ 

## That is all.

## Please enjoy Spring Break.

See you the Tuesday after, when we will talk about regular expressions.